Colourings of Path Systems

Iren Darijani and David Pike Memorial University of Newfoundland

For a graph G, a G-design of order v consists of a collection \mathcal{B} of subgraphs of the complete graph K_v such that

- each subgraph in \mathcal{B} is isomorphic to G
- the subgraphs in \mathcal{B} partition the edges of K_v

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Some families of *G*-designs:

- Steiner triple systems, denoted STS(v)
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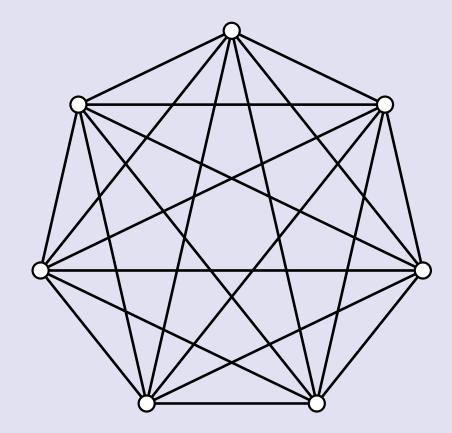
- each subgraph in $\mathcal B$ is isomorphic to G
- the subgraphs in \mathcal{B} partition the edges of K_v

Some families of *G*-designs:

- Steiner triple systems, denoted STS(v)
 G is a 3-cycle, namely C₃ or K₃
- balanced incomplete block designs, denoted BIBD(v, k, 1)*G* is a clique K_k
- path systems
 - G is a path P_m

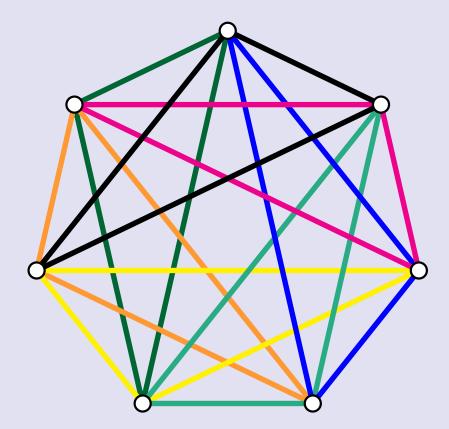
Steiner triple systems

Example: K_7

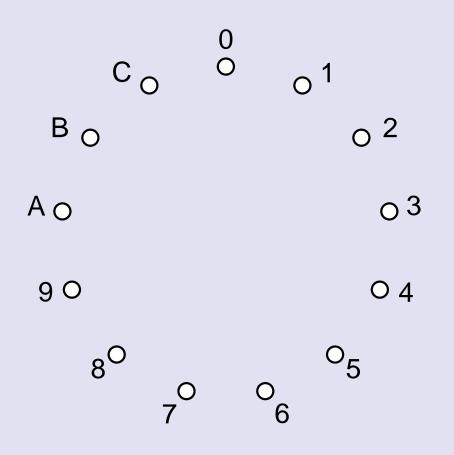


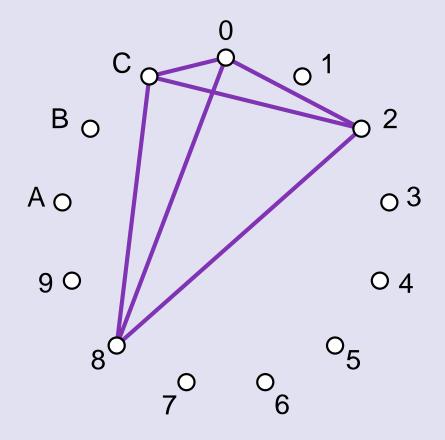
Steiner triple systems

Example: a 3-cycle decomposition of K_7

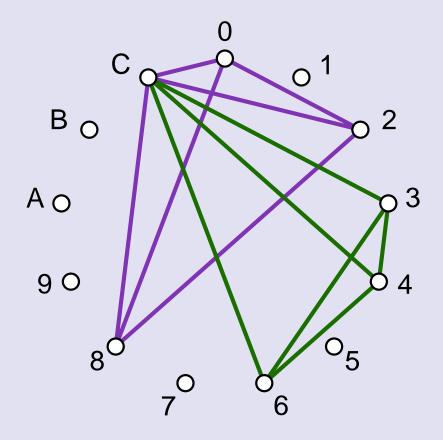


i.e., a STS(7) or a BIBD(7,3,1)



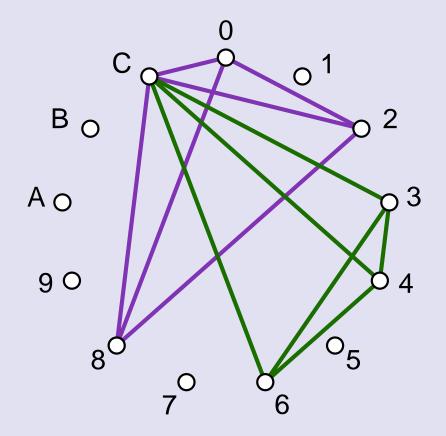


 $\{0,2,8,C\}$

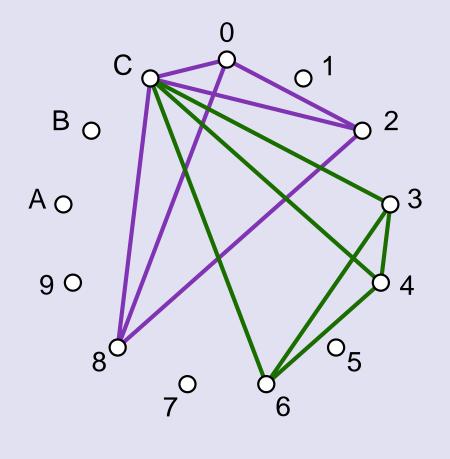


 $\{0,2,8,C\}$

 ${3,4,6,C}$

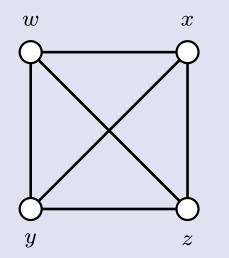


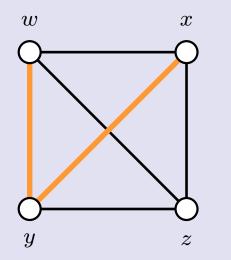
 $\{0,1,3,9\}$ {0,2,8,C} $\{0,4,5,7\}$ {0,6,A,B} {1,2,4,A} {1,5,6,8} {1,7,B,C} {2,3,5,B} {2,6,7,9} {3,4,6,C} {3,7,8,A} {4,8,9,B} {5,9,A,C}



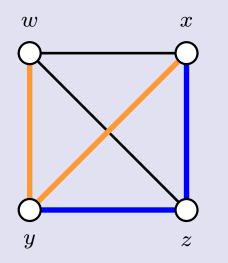
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This is equivalent to a K_4 -decomposition of K_{13}

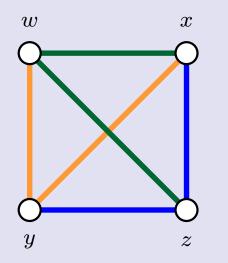




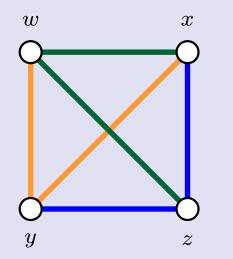
$$\mathcal{B} = \big\{ (w, y, x) \big\}$$



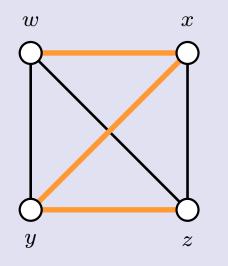
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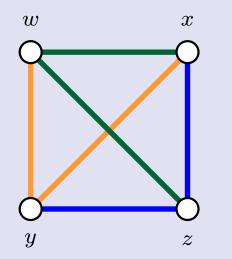
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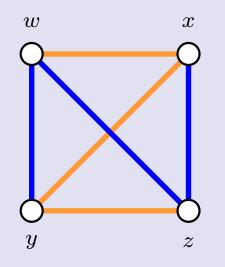
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Conditions for existence:

Theorem (Tarsi, 1983)

A P_m -decomposition of K_v exists if and only if

- $m \leqslant v$
- (m-1) divides $\binom{v}{2}$

Definition

Any integer v satisfying the above criteria will be called P_m -admissible or just admissible.

Resolvable Path Decompositions:

Theorem (Horton, 1985)

A resolvable P_3 -decomposition of K_v exists if and only if $v \equiv 9 \pmod{12}$

Theorem (Bermond, Heinrich and Yu, 1990)

A resolvable P_m -decomposition of K_v exists if and only if

- $v \equiv 0 \pmod{m}$
- $m(v-1) \equiv 0 \pmod{2(m-1)}$

A weak *c*-colouring of a design \mathcal{D} consists of a partition of the points of \mathcal{D} into *c* colour classes such that no block of \mathcal{D} is monochromatic.

A design \mathcal{D} is said to be *c*-chromatic if *c* is the smallest integer for which \mathcal{D} admits a weak *c*-colouring. Notation: we write $\chi(\mathcal{D}) = c$ if \mathcal{D} is *c*-chromatic.

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Example: $\{0,1,3,9\}$ $\{1,5,6,8\}$ $\{3,7,8,A\}$ $\{0,2,8,C\}$ $\{1,7,B,C\}$ $\{4,8,9,B\}$ $\{0,4,5,7\}$ $\{2,3,5,B\}$ $\{5,9,A,C\}$ $\{0,6,A,B\}$ $\{2,6,7,9\}$ $\{1,2,4,A\}$ $\{3,4,6,C\}$

A BIBD(13,4,1)

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This design is 2-chromatic.

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Example: $\{0,1,3\}$ $\{1,2,4\}$ $\{2,3,5\}$ $\{0,2,6\}$ $\{1,5,6\}$ $\{3,4,6\}$ $\{0,4,5\}$

A STS(7)

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Example: $\{0,1,3\}$ $\{1,2,4\}$ $\{2,3,5\}$ $\{0,2,6\}$ $\{1,5,6\}$ $\{3,4,6\}$ $\{0,4,5\}$

This design is 3-colourable. But is it 3-chromatic?

Some History – Steiner Triple Systems

• Every STS(v) with $v \ge 7$ requires at least 3 colours.

(Rosa and Pelikán, 1970)

• Every STS(v) with $7 \le v \le 15$ is 3-chromatic.

(Mathon, Phelps and Rosa, 1983)

Every STS(19) is 3-chromatic.

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For each c there is a STS with chromatic number at least c.

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• For each $c \ge 3$ there is a *c*-chromatic STS.

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- For each c there is a STS with chromatic number at least c. (Rosa, 1970)
- For each $c \ge 3$ there is a *c*-chromatic STS. (de Brandes, Phelps and Rödl, 1982)
- There is a 4-chromatic STS(21).
- (Haddad, 1999)
- There is a 5-chromatic STS(63).
- There is a 6-chromatic STS(243).

(Fugère, Haddad and Wehlau, 1994)

(Bruen, Haddad and Wehlau, 1998)

Some History – BIBDs

- For each admissible v, i.e. $v \equiv 1$ or 4 (mod 12), there is a 2-chromatic BIBD(v, 4, 1). (Hoffman, Lindner and Phelps, 1990) (Franek, Griggs, Lindner and Rosa, 2002)
- A 3-chromatic BIBD(v, 4, 1) exists if and only if $v \equiv 1 \text{ or } 4 \pmod{12}$ and $v \ge 25$. (Rodger, Wantland, Chen, Zhu, 1994)
- The obvious necessary conditions for the existence of a ${\sf BIBD}(v,4,1)$ are asymptotically sufficient for the existence of a c-chromatic ${\sf BIBD}(v,4,1)$. (Linek and Wantland, 1998)
- For each admissible v, i.e. $v \equiv 1$ or 5 (mod 20), there is a 2-chromatic BIBD(v, 5, 1). (Ling, 1999)

Some History – BIBDs

- For $\lambda \ge 2$, for each admissible v, (Hoffman, Lindner and Phelps, 1990) there is a 2-chromatic BIBD $(v, 4, \lambda)$. (Hoffman, Lindner and Phelps, 1991) (Rosa and Colbourn, 1992)
- For all integers $\lambda \ge 1$, $c \ge 2$ and $k \ge 3$ with $(c, k) \ne (2, 3)$, the obvious necessary conditions for the existence of a BIBD (v, k, λ) are asymptotically sufficient for the existence of a *c*-chromatic BIBD (v, k, λ) . (Horsley and Pike, 2014)

Path Systems

Theorem (Darijani and Pike, 2023+)

For each $c \ge 2$ and $m \ge 3$ there exists a *c*-chromatic P_m system.

Path Systems

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For each $c \ge 2$ and $m \ge 3$ there exists a *c*-chromatic P_m system.

Proof Sketch for when m is even:

For sufficiently large v, there exists a c-chromatic BIBD(v, m, 1) (by Horsley and Pike, 2014).

The graph K_{m+1} can be decomposed into Hamilton cycles (by Walecki, 1890s). Deleting one vertex from K_{m+1} yields a decomposition of K_m into P_m paths.

By decomposing each block of the BIBD(v, m, 1) into P_m paths, we obtain a *c*-chromatic P_m system of order *v*.

Path Systems

Theorem (Darijani and Pike, 2023+)

For each $c \ge 2$ and $m \ge 3$ there exists a *c*-chromatic P_m system.

Proof Sketch for when m is odd and $m \ge 5$:

For sufficiently large v, there exists a c-chromatic BIBD(v, m, 1) such that $v \equiv 0 \pmod{m}$ and $v - 1 \equiv 0 \pmod{2m - 2}$.

The block set of this BIBD can be partitioned into pairs of blocks B and B' that share one point.

The graph K_m can be decomposed into Hamilton cycles. By removing a particular edge from each Hamilton cycle, a set of P_m paths is obtained from B and B'.

We obtain a *c*-chromatic P_m system of order *v*.

P_4 Path Systems

Theorem (Darijani and Pike, 2023+)

For each $c \ge 2$, there exists a *c*-chromatic P_4 system of order v for all sufficiently large admissible v.

Recall: The admissible orders for a P_4 system are $v \equiv 0, 1 \pmod{3}$

Proof Sketch:

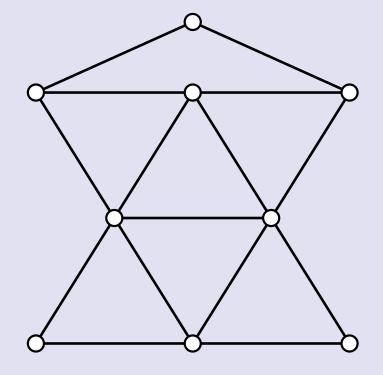
For c = 2, we show that there exist 2-chromatic systems for all admissible orders (by direct construction).

For $c \ge 3$, we show how to use a *c*-chromatic system of order *v* to construct *c*-chromatic systems for some of the next admissible orders. We then iterate this process, beginning with initial instances (such as those from the previous slides).

A *c*-colouring of a design is unique if every *c*-colouring of the design has the same partition into colour classes.

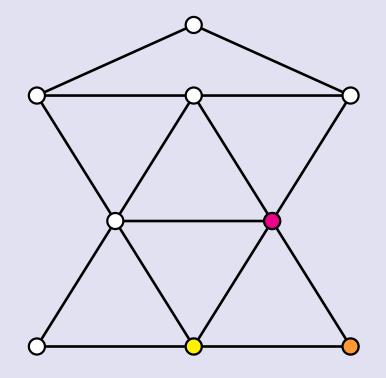
Note that whenever a design is uniquely c-colourable, it must also be c-chromatic.

An easy example of a uniquely 3-colourable graph

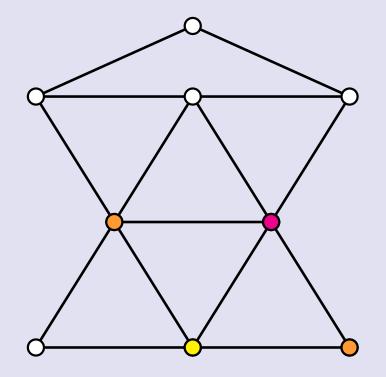


An easy example of a uniquely 3-colourable graph

At least 3 colours are necessary for a proper colouring.

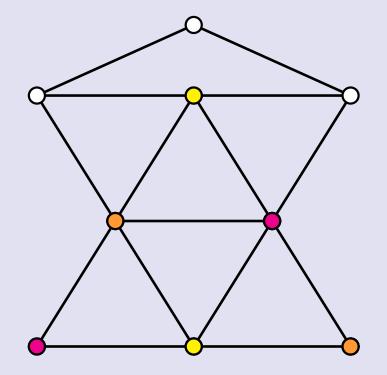


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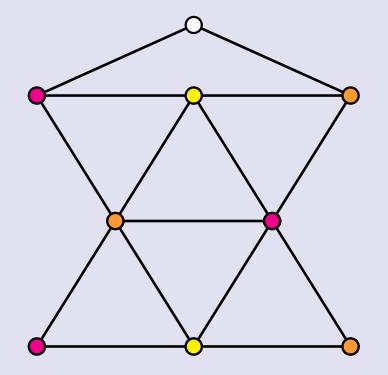
We now have no choice.

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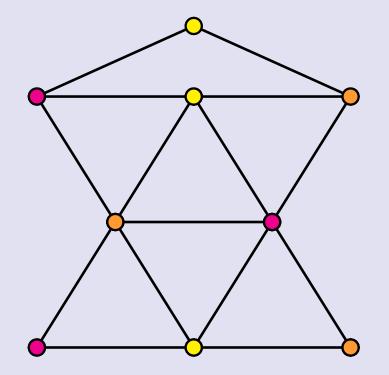
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This is a unique partition into 3 colour classes.

Some History – Unique Colourings of Designs

• A uniquely 3-colourable STS(33) is observed to exist.

(de Brandes, Phelps and Rödl, 1982)

(Colbourn, Haddad and Linek, 1997)

• For each $v \ge 25$ there exists a uniquely 3-colourable STS(v) (Forbes, 2003)

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As far as I am aware, when $c \ge 4$, no examples of uniquely *c*-colourable STS are currently known to exist.

Uniquely *c*-colourable BIBD(v, k, 1) with $k \ge 4$ are also unknown.

Uniquely *c*-colourable *e*-star systems were studied.

(Darijani and Pike, 2020)

P₄ Path Systems

Theorem (Darijani and Pike, 2023+)

There exists a uniquely 2-colourable P_4 system of order v for each admissible $v \ge 109$.

Uniquely 2-Colourable *P*₄ Systems

- We built a system S_{28} of order 28 in which two specific points cannot have the same colour in any 2-colouring.
- We then used S₂₈ to build a partial system on 109 points that has a unique 2-colouring.
- We successfully completed this partial system to a full system of order 109, by adding more blocks, none of which are monochromatic.
- Finally we showed how to iteratively take a uniquely colourable system and embed it in slightly larger systems that are also uniquely colourable.

Building S_{28}

- Suppose points 27 and 28 have the same colour, say white.
- For i = 1, 2, ..., 12, add the block (27, 2i, 2i 1, 28). This forces one of points 1 & 2 to be black, one of points 3 & 4 to be black, etc.
- For distinct $i, j, \ell \in \{1, 2, \dots, 12\}$ consider these blocks:

 $(2i - 1, 2j, 2\ell, 2i) \qquad (2i - 1, 2j - 1, 2\ell - 1, 2i)$ $(2i - 1, 2\ell - 1, 2j, 2i) \qquad (2i - 1, 2\ell, 2j - 1, 2i)$

Adding these blocks for $(i, j, \ell) = (1, 2, 3)$ forces points 1 & 2 to have different colours.

Do similar for $(i, j, \ell) = (4, 5, 6), (7, 8, 9), (10, 11, 12)$ so that points 7 & 8, 13 & 14, 19 & 20 are different too.

Building S_{28} (continued)

- Specifying some more blocks forces points 1,7,13,19 to be one colour and 2,8,14,20 to be the other.
- Similar steps force points 3,9,15,21 to be one colour and 4,10,16,22 to be the other.
- A monochromatic block is eventually forced.
 Hence the supposition that 27 & 28 have the same colour is contradicted.
- Demonstrating a valid 2-colouring where 27 and 28 are different is all that remains, to confirm that the system's chromatic number is not larger than 2. In fact, we exhibit a 2-colouring for which the colour classes have equal size.

Building S_{109}

- Take four copies of S_{28} on point sets $\{1, 2, \dots, 27, 28\}$ $\{1', 2', \dots, 27', 28\}$ $\{1'', 2'', \dots, 27'', 28\}$ $\{1''', 2''', \dots, 27''', 28\}$
- If point 28 is white, then 27, 27', 27'' and 27''' are each black.
- Add blocks between the four copies of S_{28} such as:

 $(27, 1', 27'', 27') \qquad (27, 3', 27'', 27''')$ $(27, 1'', 27''', 27') \qquad (27', 27, 1''', 27'')$

Points 1', 3', 1'' and 1''' are therefore white.

- Add more blocks, eventually forcing each point's colour.
- Then complete the design by adding additional blocks (none of which are monochromatic).

Future Work

- Prove that uniquely c-colourable P_4 systems exist for $c \ge 3$.
- Also consider P_m systems with $m \neq 4$

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- Prove that uniquely c-colourable P_4 systems exist for $c \ge 3$.
- Also consider P_m systems with $m \neq 4$
- Prove that uniquely *c*-colourable Steiner triple systems exist for *c* ≥ 4.
- Also consider (v, k, λ) -BIBDs with $k \ge 4$ and *m*-cycle systems with $m \ge 4$.

Thank you.

Acknowledgements:



